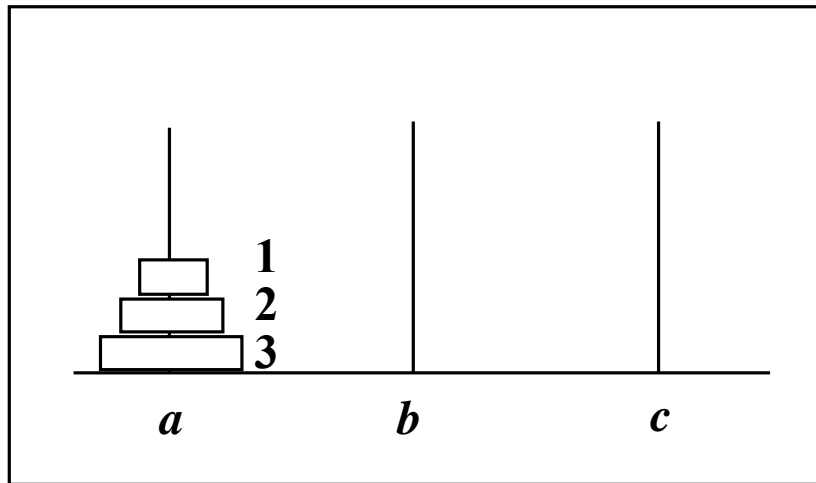


Tower of Hanoi Problem (The general case)

The problem is as follows. There are three pivots a , b , and c . On the first one there is a set of n disks, each of different radius. The task is to move all the disks to the pivot c moving only one disk at a time. In addition, at no time during the process may a disk be placed on top of a smaller disk. The pivot c can, of course, be used as a temporary resting place for the disks.



Let us designate an elementary step of moving disk number i from the pivot x to the pivot y as $p(i, x, y)$, a terminal symbol with parameters. Thus a solution of the Tower of Hanoi Problem might be represented as the following string of symbols with parameters:

$$p(i_1, x_1, y_1)p(i_2, x_2, y_2)\dots p(i_m, x_m, y_m).$$

This is the string of the language of all possible sequences of moves. Consider the controlled grammar shown in Figure 4. We will apply this grammar for derivation of a solution for the case of three disks: $n=3$, $x=a$, $y=c$. It means that the values of parameters for the starting symbol S are $S(3, a, b)$.

Controlled grammar generating solutions to the Tower of Hanoi Problem

L	Q	Kernel, π_k	π_n	F_T	F_F
1	Q_1	$S(n, x, y) \rightarrow A(n, x, y)$		2	\emptyset
2	Q_2	$A(n, x, y) \rightarrow A(f_1(n), x, f_2(x, y))$ $p(n, x, y)$ $A(f_1(n), f_2(x, y), y)$		2	3
3	Q_3	$A(n, x, y) \rightarrow p(n, x, y)$		2	\emptyset

Here $V_T = \{p\}$

$V_N = \{S, A\}$

V_{PR}

$Pred = \{Q_1, Q_2, Q_3\}$,

$Q_1 = T$

$Q_2(n) = T$, if $n > 1$; $Q_2(n) = F$, if $n = 1$.

$Q_3(n) = T$, if $n = 1$; $Q_3(n) = F$, if $n > 1$.

$Var = \{n, x, y\}$

$F = Fcon \cup Fvar$,

$Fcon = \{f_1, f_2\}$

$f_1(n) = n-1, n = 2, 3, \dots$

$f_2(x, y)$ yields the value from $\{a, b, c\} \setminus \{x, y\}$, where values of

x, y are from $\{a, b, c\}$

$Fvar = \{3, a, c\}$

$E = \mathbf{Z}_+ \cup \{a, b, c\}$

Parm: $S \rightarrow Var, A \rightarrow Var, p \rightarrow Var$

$L = \{1, 2, 3\}$

At the beginning of derivation: $x = a, y = c, n = 3$.

Generation of a solution in case of n = 3:

$$S(3, a, c) \stackrel{1}{\Rightarrow} A(3, a, c) \stackrel{2}{\Rightarrow} A(2, a, b)p(3, a, c)A(2, b, c)$$

$$\stackrel{2}{\Rightarrow} A(1, a, c)p(2, a, b)A(1, c, b)p(3, a, c)A(2, b, c)$$

$$\stackrel{3}{\Rightarrow} p(1, a, c)p(2, a, b)A(1, c, b)p(3, a, c)A(2, b, c)$$

$$\stackrel{3}{\Rightarrow} p(1, a, c)p(2, a, b)p(1, c, b)p(3, a, c)A(2, b, c)$$

$$\stackrel{2}{\Rightarrow} p(1, a, c)p(2, a, b)p(1, c, b)p(3, a, c)A(1, b, a)p(2, b, c)A(1, a, c)$$

$$\stackrel{3}{\Rightarrow} p(1, a, c)p(2, a, b)p(1, c, b)p(3, a, c)p(1, b, a)p(2, b, c)A(1, a, c)$$

$$\stackrel{3}{\Rightarrow} p(1, a, c)p(2, a, b)p(1, c, b)p(3, a, c)p(1, b, a)p(2, b, c)p(1, a, c).$$

Generation of a solution in case of n = 4:

$$S(4, a, c) \stackrel{1}{=} A(4, a, c) \stackrel{2}{=} \underline{A(3, a, b)} p(4, a, c) A(3, b, c)$$

$$\begin{aligned} \underline{A(3, a, b)} &\stackrel{2}{=} A(2, a, c) p(3, a, b) A(2, c, b) \\ &\stackrel{2}{=} A(1, a, b) p(2, a, c) A(1, b, c) p(3, a, b) A(2, c, b) \\ &\stackrel{3}{=} p(1, a, b) p(2, a, c) A(1, b, c) p(3, a, b) A(2, c, b) \\ &\stackrel{3}{=} p(1, a, b) p(2, a, c) p(1, b, c) p(3, a, b) A(2, c, b) \\ &\stackrel{2}{=} p(1, a, b) p(2, a, c) p(1, b, c) p(3, a, b) A(1, c, a) p(2, c, b) A(1, a, b) \\ &\stackrel{3}{=} p(1, a, b) p(2, a, c) p(1, b, c) p(3, a, b) p(1, c, a) p(2, c, b) A(1, a, b) \\ &\stackrel{3}{=} p(1, a, b) p(2, a, c) p(1, b, c) p(3, a, b) p(1, c, a) p(2, c, b) p(1, a, b). \end{aligned}$$

$$\begin{aligned} 2, \dots, 3 &\Rightarrow p(1, a, b) p(2, a, c) p(1, b, c) p(3, a, b) p(1, c, a) p(2, c, b) p(1, a, b) \\ &\quad p(4, a, c) \underline{A(3, b, c)} \end{aligned}$$

$$\begin{aligned} \underline{A(3, b, c)} &\stackrel{2}{=} A(2, b, a) p(3, b, c) A(2, a, c) \\ &\stackrel{2}{=} A(1, b, c) p(2, b, a) A(1, c, a) p(3, b, c) A(2, a, c) \\ &\stackrel{3}{=} p(1, b, c) p(2, b, a) A(1, c, a) p(3, b, c) A(2, a, c) \\ &\stackrel{3}{=} p(1, b, c) p(2, b, a) p(1, c, a) p(3, b, c) A(2, a, c) \\ &\stackrel{2}{=} p(1, b, c) p(2, b, a) p(1, c, a) p(3, b, c) A(1, a, b) p(2, a, c) A(1, b, c) \\ &\stackrel{3}{=} p(1, b, c) p(2, b, a) p(1, c, a) p(3, b, c) p(1, a, b) p(2, a, c) A(1, b, c) \\ &\stackrel{3}{=} p(1, b, c) p(2, b, a) p(1, c, a) p(3, b, c) p(1, a, b) p(2, a, c) p(1, b, c). \end{aligned}$$

$$\begin{aligned} 2, \dots, 3 &\Rightarrow p(1, a, b) p(2, a, c) p(1, b, c) p(3, a, b) p(1, c, a) p(2, c, b) p(1, a, b) \\ &\quad p(4, a, c) \\ &\quad p(1, b, c) p(2, b, a) p(1, c, a) p(3, b, c) p(1, a, b) p(2, a, c) p(1, b, c). \end{aligned}$$

Induction

Let's prove by induction that $1+2+3+\dots+n = n(n+1)/2$

1. For $n=1$ this is true: $1=1(1+1)/2$

2. Assume that the statement is true for $n=m$, i.e.,
 $1+2+3+\dots+m = m(m+1)/2$

3. Let's prove that the statement is true for $n=m+1$.

Indeed, from the assumption 2 we have

$$1+2+3+\dots+m + (m+1) = m(m+1)/2 + (m+1) = \\ (m/2+1)(m+1) = [(m+2)/2](m+1) = (m+1)(m+2)/2.$$

The statement is proved for $n = m+1$.

4. By induction we conclude that this statement is true for all $n = 1, 2, 3, \dots$

Consider the following problem:

$$1+3+5+7+\dots+(2n-1) = ?$$

Find the answer and prove by induction.

Solution

1. $1 = 1^2$

2. Assume that $1+3+5+\dots+(2m-1) = m^2$

3. $1+2+3+\dots+(2(m+1)-1) = 1+2+3+\dots+(2m-1) + (2m+1) = m^2 + (2m+1) = \\ m^2 + 2m + 1 = (m+1)^2$

Examples of Induction

Prove that $1^2 + 2^2 + \dots + n^2 = [n(n+1)(2n+1)] / 6$

Prove that $1^3 + 2^3 + \dots + n^3 = [n^2(n+1)^2] / 4$

Prove that the number that consists of 3^n units (like $\frac{1111\dots1}{3^n}$) is divisible by 3^n .

Solution

1. 111 is divisible by 3^1 .
2. Assume that the statement is true for $n = m$, i.e., $A = \frac{111 \dots 1}{3^m}$.
3. Let's prove the statement for $B = \frac{111 \dots 1}{3^{m+1}}$.

Consider

$$\frac{111111111}{3^{1+1}} = \frac{111000000}{3^1} + \frac{111000}{2 \cdot 3^1} + \frac{111}{3^1} + \frac{111}{3^1} + \frac{111}{3^1}$$

Let $a = 111$, then

$$\begin{aligned} &= a \cdot 103^1 \cdot 2 + a \cdot 103^1 + a \\ &= a \times (103^1 \cdot 2 + 103^1 + 1) \end{aligned}$$

$$B = \frac{111 \dots 1000 \dots 0}{3^m} + \frac{111 \dots 1000 \dots 0}{2 \cdot 3^m} + \frac{111 \dots 1}{3^m} + \frac{111 \dots 1}{3^m} + \frac{111 \dots 1}{3^m} = A \cdot 103^{m \cdot 2} + A \cdot 103^m + A$$

By assumption A is divisible by 3^m ; hence, $B = A \times (103^{m \cdot 2} + 103^m + 1)$ is divisible by 3^{m+1} , because the sum of digits of the expression in parenthesis is equal to 3, i.e., it is divisible by 3.

Induction for the Tower of Hanoi problem

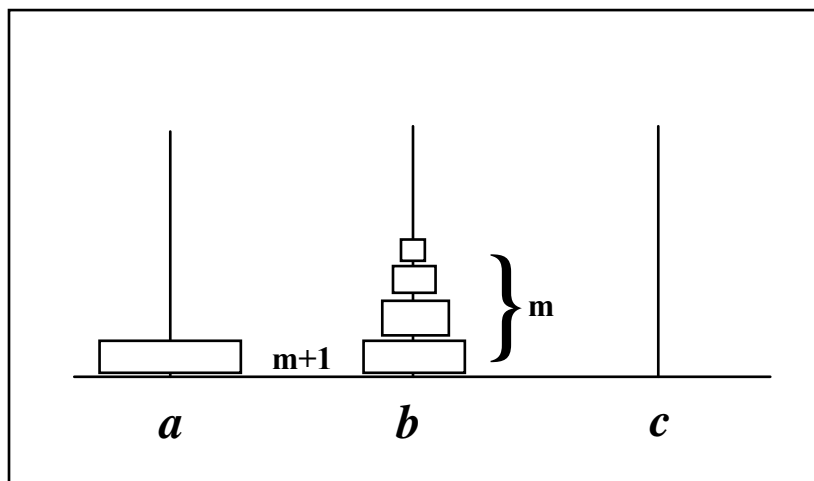
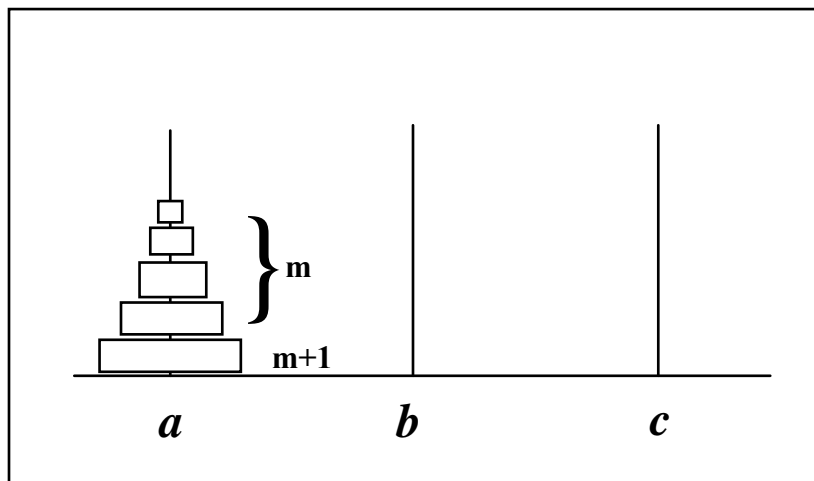
Let's prove that the grammar generates a solution of this problem in general case. We shall prove by induction.

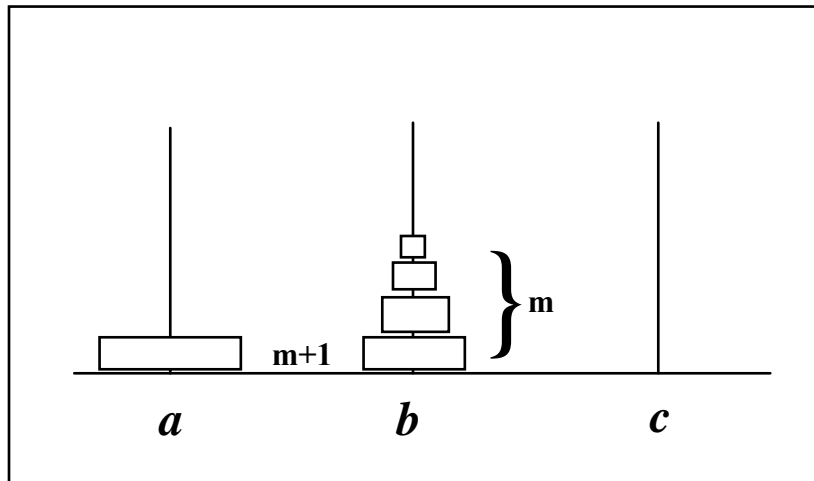
1. We proved that the grammar generates a solution of this problem for $n = 3$.
2. Assume that the grammar generates a solution for $n = m$.
3. Let's prove that it generates a solution for $n = m + 1$

Consider the derivation in case of $n = m + 1$:

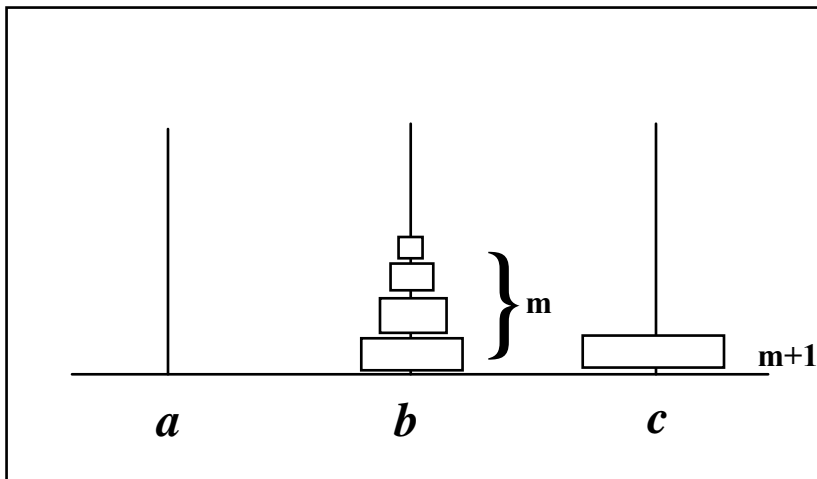
$$S(m+1, a, c) \stackrel{1}{\Rightarrow} A(m+1, a, c) \stackrel{2}{\Rightarrow} A(m, a, b)p(m+1, a, c)A(m, b, c).$$

Obviously, the following application of the grammar to the symbol $A(m, a, b)$ will generate the string of symbols. According to the assumption of induction 2. this string corresponds to the solution of the Tower of Hanoi problem with m disks on the pivot a . These disks must be moved to the pivot b .



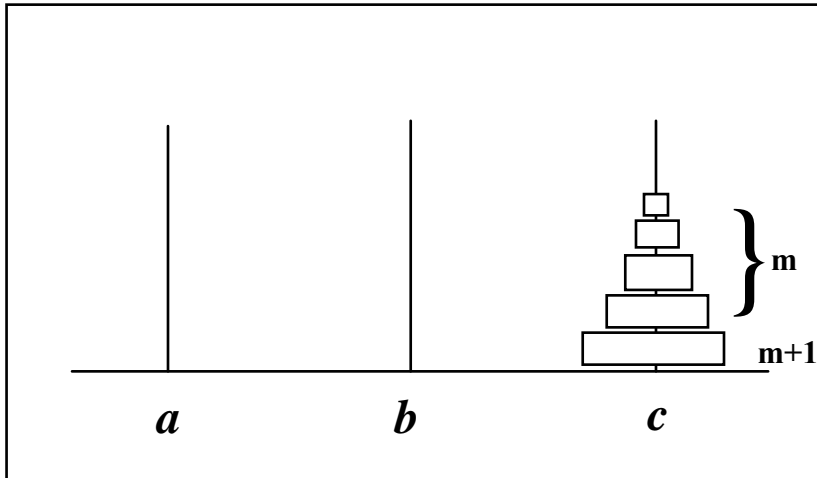
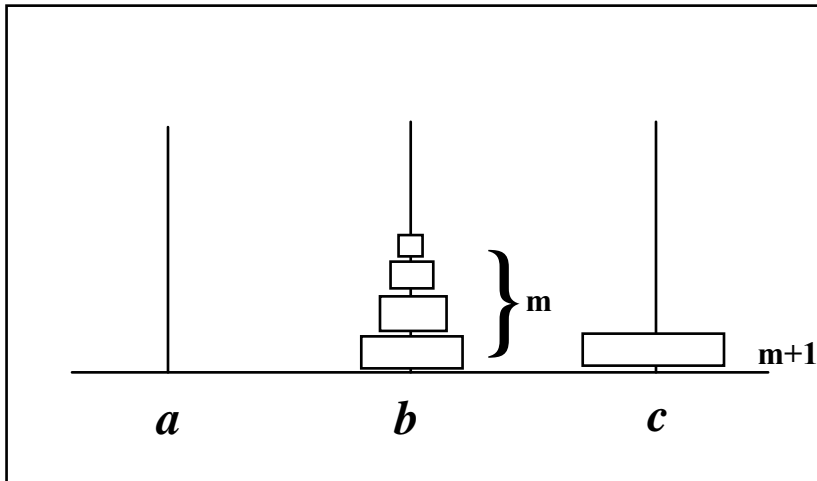


When these disks are moved to the pivot b we can apply $p(m+1, a, c)$, i.e., we can move disk $m+1$ from pivot a to pivot c .



$$S(m+1, a, c) \stackrel{1}{\Rightarrow} A(m+1, a, c) \stackrel{2}{\Rightarrow} A(m, a, b)p(m+1, a, c)A(m, b, c).$$

The following application of the grammar to the symbol $A(m, b, c)$ will generate the string of symbols. According to the assumption of induction 2. this string corresponds to the solution of the Tower of Hanoi problem with m disks on the pivot b . These disks must be moved to the pivot c .



It means that the grammar generates a solution for $n = m+1$.

4. Conclusion: by induction the grammar generates a solution of the Tower of Hanoi problem for all $n = 3, 4, 5, \dots$